



Chapter 4

Testability Analysis

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Purpose

- Measurements used to guide
 - Automatic Test Pattern Generation (ATPG)
 - Design-for-Testability (DfT)
- Testability?
 - Characteristic of a circuit that influences various costs associated with the test (length, complexity of generating a sequence, etc.)



Measurements

- Controllability

- Indicates the relative difficulty of positioning a line at 0 or at 1 from the Primary Inputs (PIs)

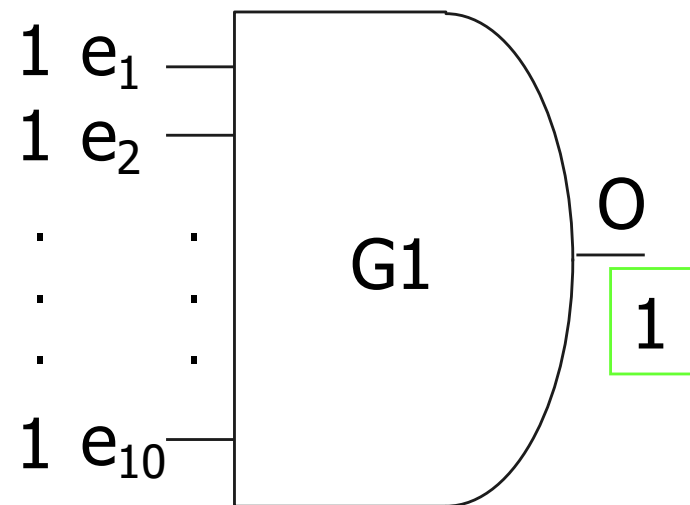
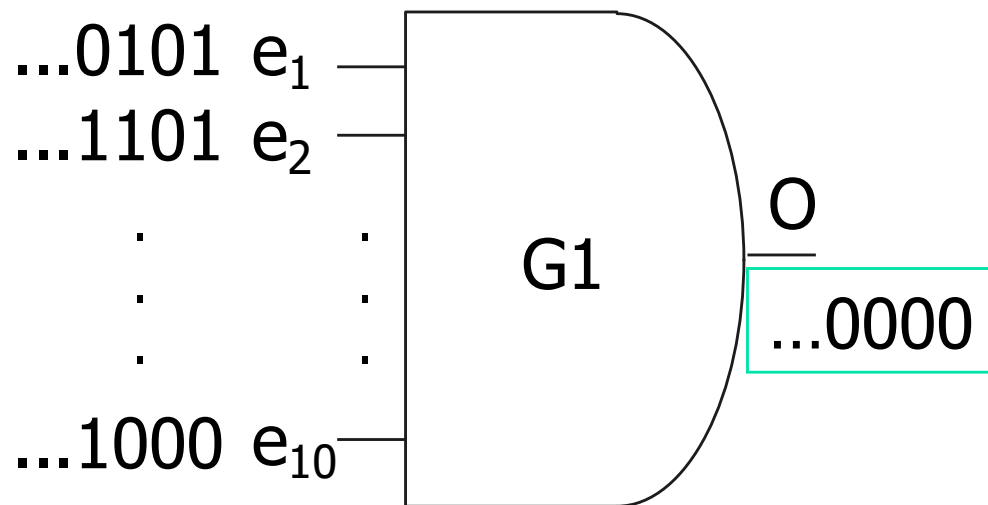
- Observability

- Indicates the relative difficulty of propagating an error from a line to the Primary Outputs (POs)

Example

Sa0 test at a AND-gate output

- Generation of random vectors
 - Control-at-1 at G1 output is "hard" ($1/2^n$ probability)
- Generation of deterministic vectors
 - Control-at-1 at G1 output is "simple"





Computation Methods

- Testability measurements for a deterministic test
 - **SCOAP** (Sandia Controllability and Observability Analysis Program)
- Testability measurements for a random or pseudo-random test
 - **COP** (Controllability Observability Probability)



SCOAP Principle

- Based on a structural analysis of the DUT
- DUT = cells (combinatorial, sequential)
+ nets (combinatorial, sequential)
- Increasing SCOP measurements with the required testing effort



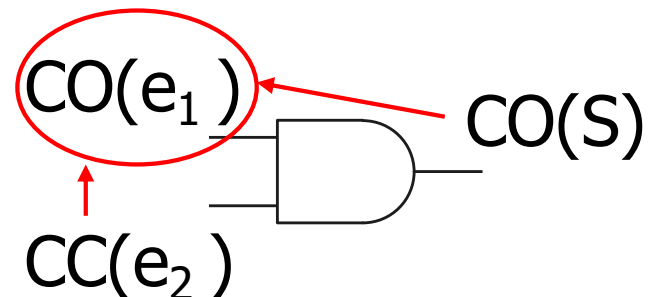
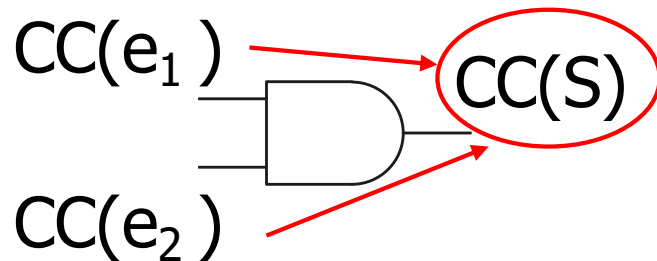
SCOAP Measurements

- Six measurements associated to each net:
 - $CC0(N)$ = Combinatorial Controllability at 0 = # minimum of comb. nets that must be controlled to bring a logic 0 at net N
 - $CC1(N)$ = Combinatorial Controllability at 1 = # minimum of comb. nets that must be controlled to bring a logic 1 at net N
 - $CO(N)$ = Combinatorial Observability = # minimum of comb. nets that must be controlled so that the effect of the fault on net N is propagated towards a PO

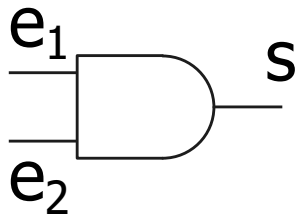
 - $SC0(N)$ = Sequential Controllability at 0 = # minimum of seq. nets that must be controlled to bring a logic 0 at net N
 - $SC1(N)$ = Sequential Controllability at 1 = # minimum of seq. nets that must be controlled to bring a logic 1 at net N
 - $SO(N)$ = Sequential Observability = # minimum of seq. nets that must be controlled so that the effect of the fault on net N is propagated towards a PO

SCOAP Evaluation

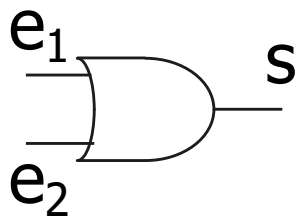
- Controllability at the output of a cell as a function of the controllability of its inputs
- Observability at the input of a cell as a function of the observability at the output and the controllability of the other inputs



SCOAP Combinatorial Cells



- $CC0(s) = \min (CC0(e1),CC0(e2)) + 1$
- $CC1(s) = CC1(e1) + CC1(e2) + 1$
- $CO(e1) = CC1(e2) + CO (s) + 1$



- $CC0(s) = CC0 (e1) + CC0 (e2) + 1$
- $CC1(s) = \min (CC1(e1),CC1(e2))+1$
- $CO(e1) = CC0(e2) + CO (s) + 1$



Classical Logic Gates

Combinatorial
Controllability at 0

AND2	$CC0(OUT) = \min(CC0(IN1), CC0(IN2)) + 1$
NAND2	$CC0(OUT) = CC1(IN1) + CC1(IN2) + 1$
OR2	$CC0(OUT) = CC0(IN1) + CC0(IN2) + 1$
NOR2	$CC0(OUT) = \min(CC1(IN1), CC1(IN2)) + 1$
INV	$CC0(OUT) = CC1(IN) + 1$
BUF	$CC0(OUT) = CC0(IN) + 1$

Combinatorial
Controllability at 1

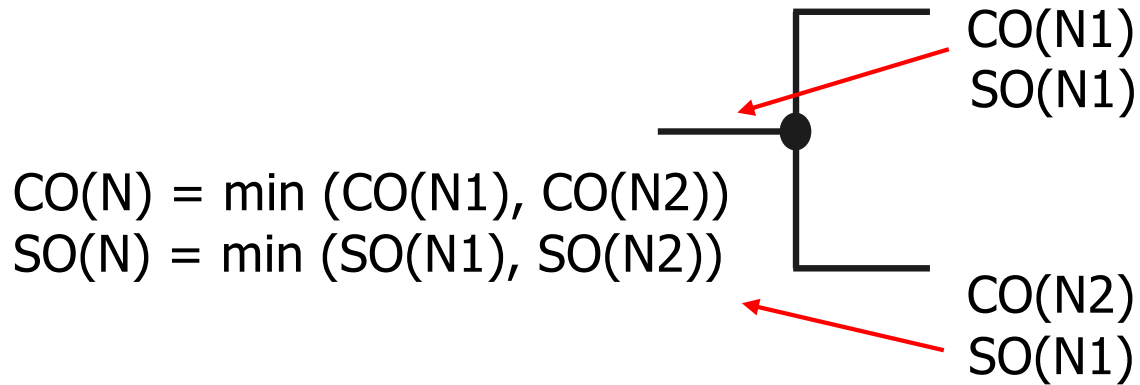
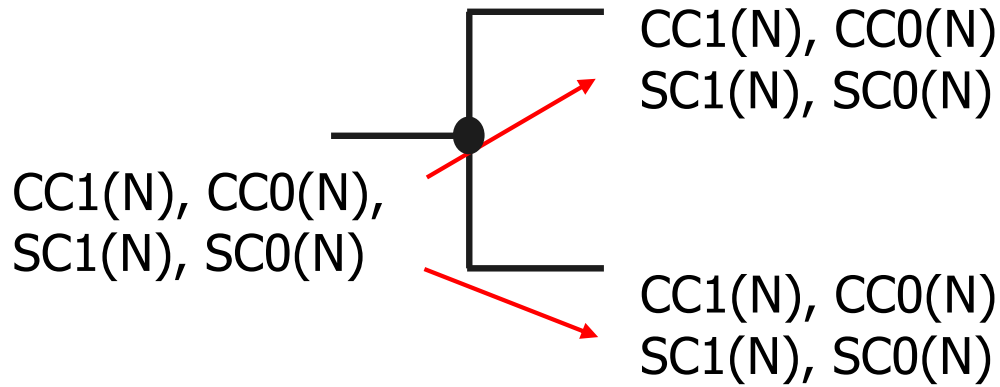
AND2	$CC1(OUT) = CC1(IN1) + CC1(IN2) + 1$
NAND2	$CC1(OUT) = \min(CC0(IN1), CC0(IN2)) + 1$
OR2	$CC1(OUT) = \min(CC1(IN1), CC1(IN2)) + 1$
NOR2	$CC1(OUT) = CC0(IN1) + CC0(IN2) + 1$
INV	$CC1(OUT) = CC0(IN) + 1$
BUF	$CC1(OUT) = CC1(IN) + 1$

Combinatorial
Observability

AND2	$CO(IN1) = CC1(IN2) + CO(OUT) + 1$
NAND2	$CO(IN1) = CC1(IN2) + CO(OUT) + 1$
OR2	$CO(IN1) = CC0(IN2) + CO(OUT) + 1$
NOR2	$CO(IN1) = CC0(IN2) + CO(OUT) + 1$
INV	$CO(IN1) = CO(OUT) + 1$
BUF	$CO(IN1) = CO(OUT) + 1$



SCOAP Fanout





SCOAP Process

- Net initialisation:

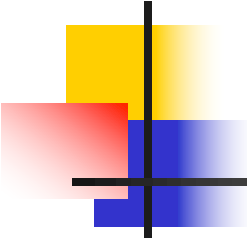
- If $N = \text{PI}$ → $\text{CC0}(N) = \text{CC1}(N) = 1$
→ $\text{SC0}(N) = \text{SC1}(N) = 0$
- If $N = \text{PO}$ → $\text{CO}(N) = \text{SO}(N) = 0$
- Else → $\text{CC0} = \text{CC1} = \text{SC0} = \text{SC1} = \infty$

- Phase 1:

- Controllability computation from PIs to POs

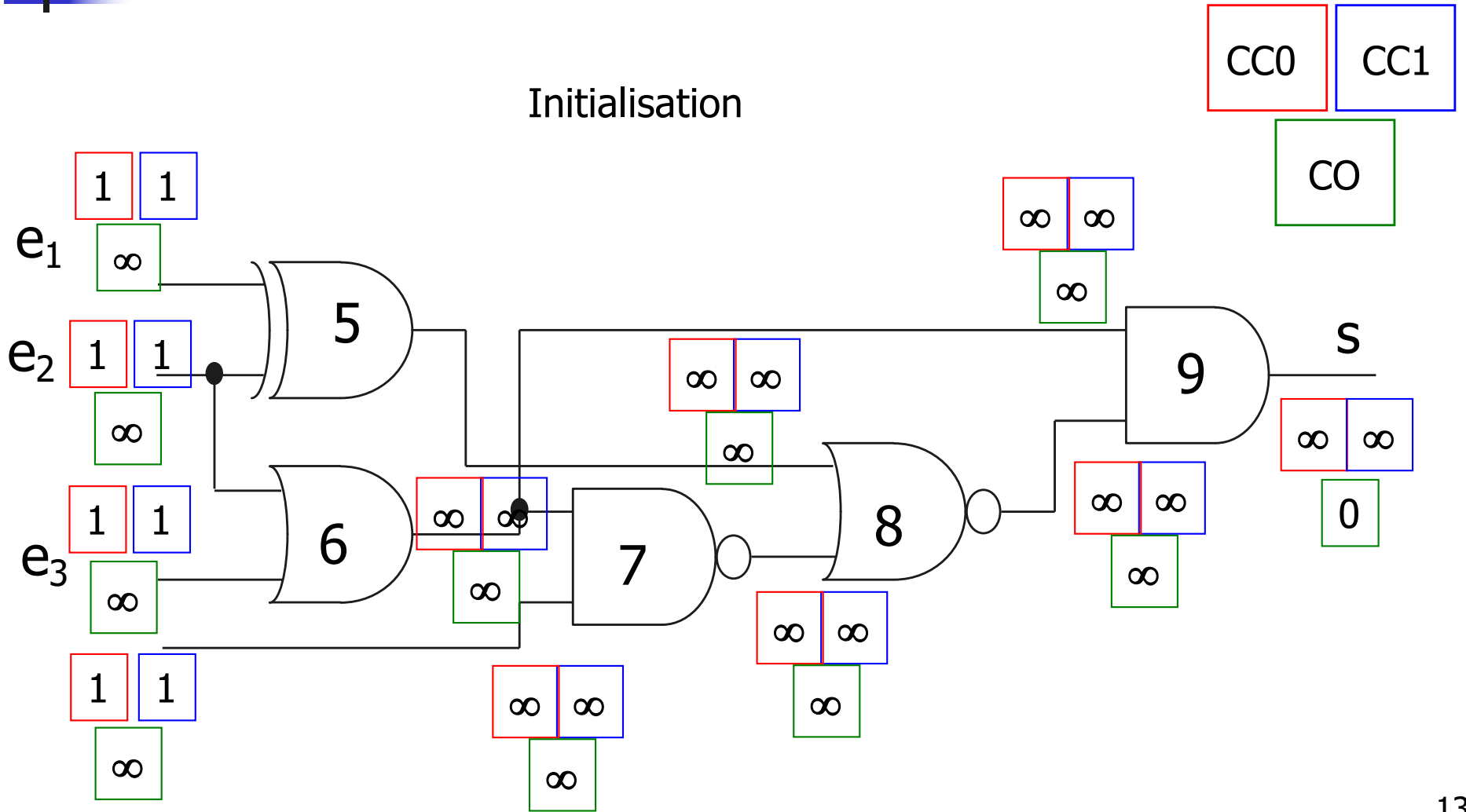
- Phase 2:

- Observability computation from POs to PIs

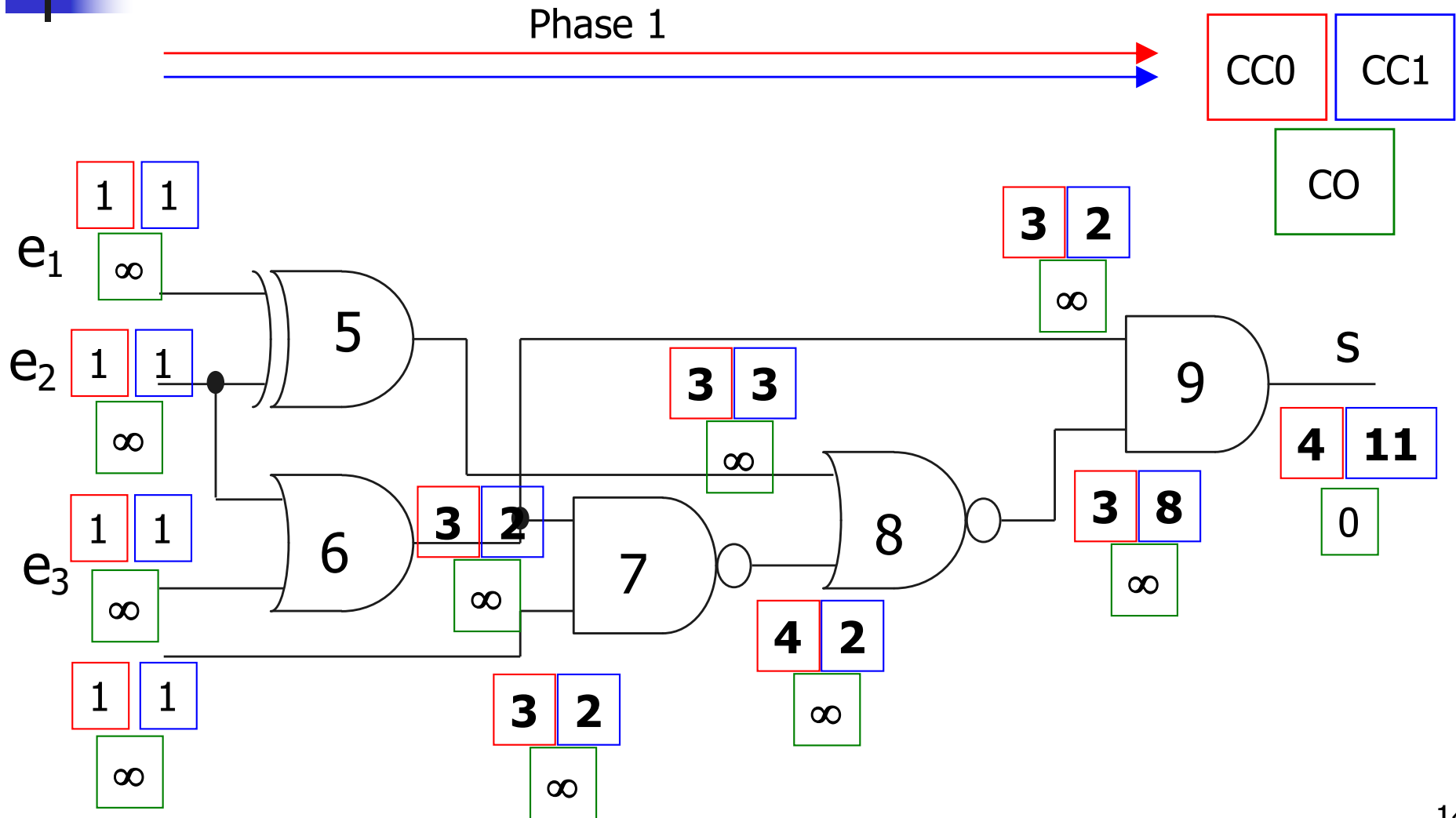


Example

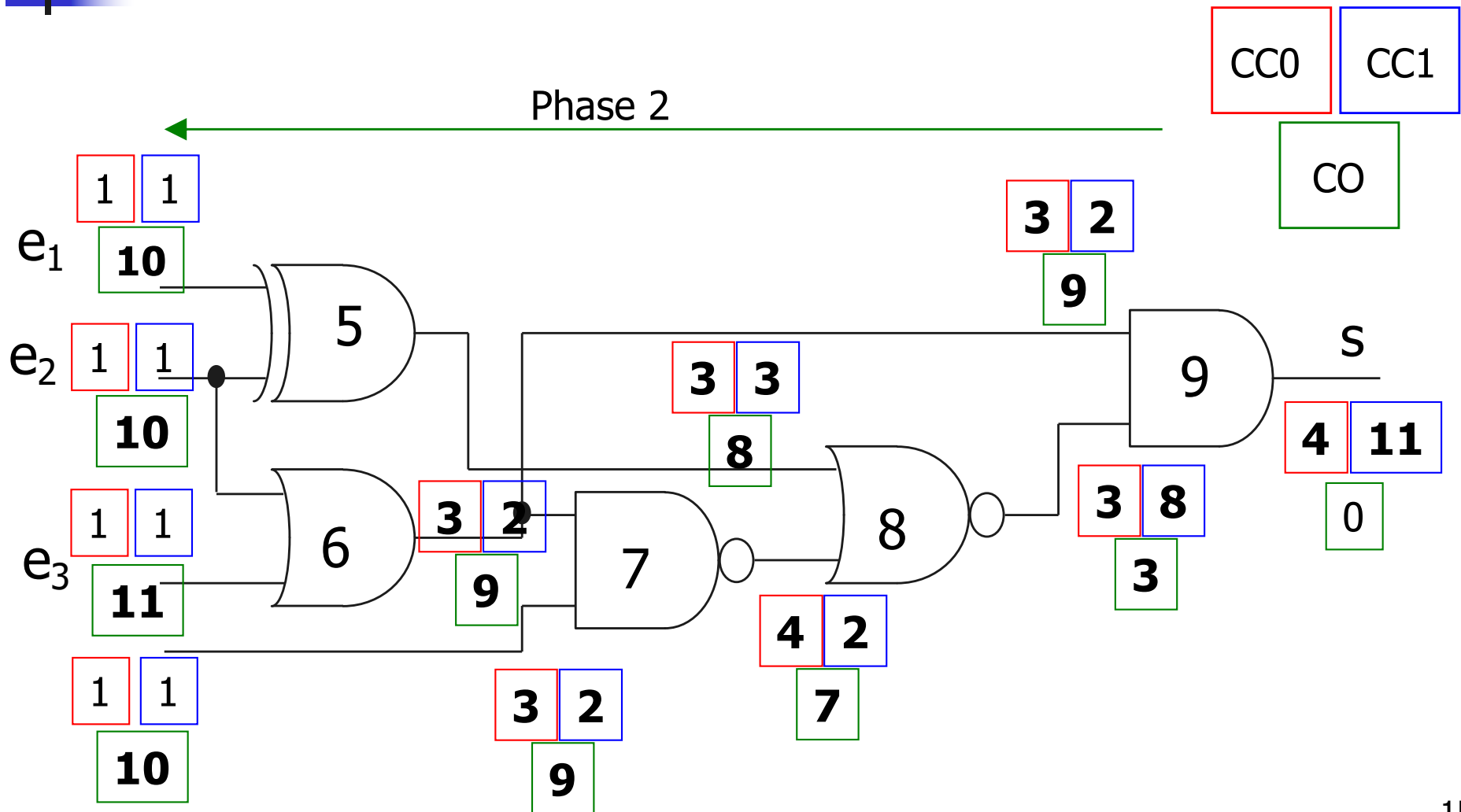
Initialisation



Example

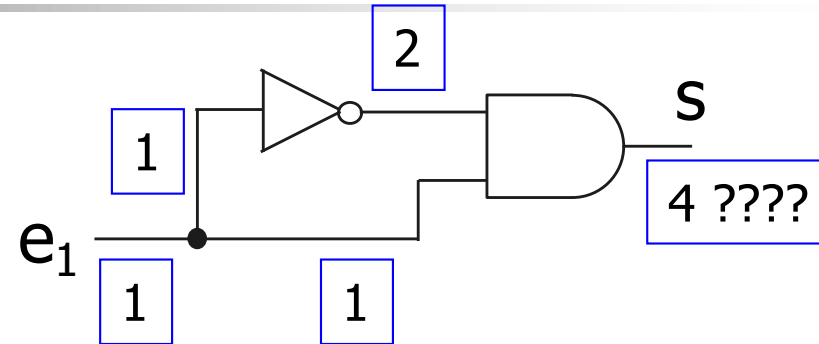


Example

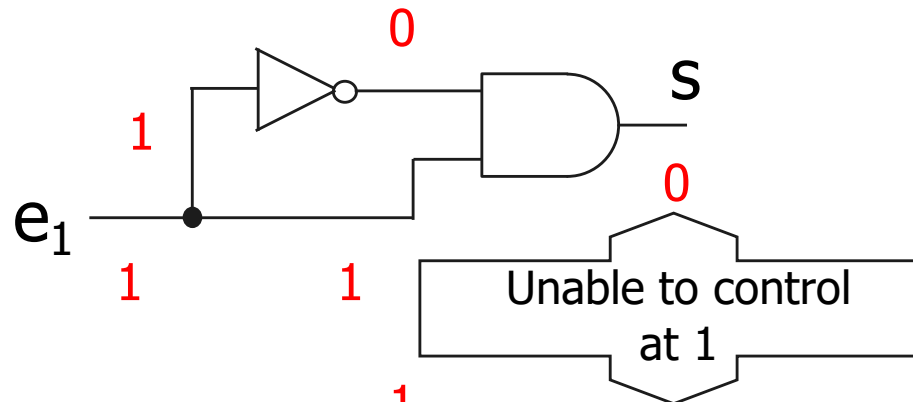


SCOAP Pros and Cons (1)

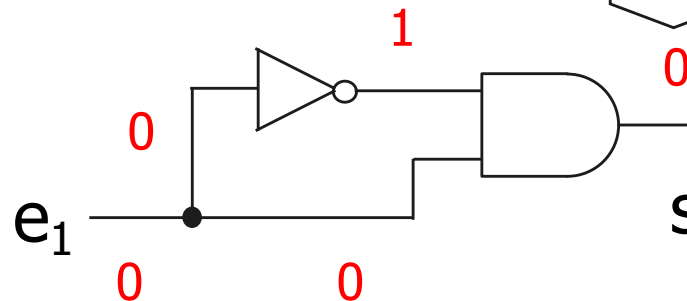
SCOAP CC1



$e_1 = 1$



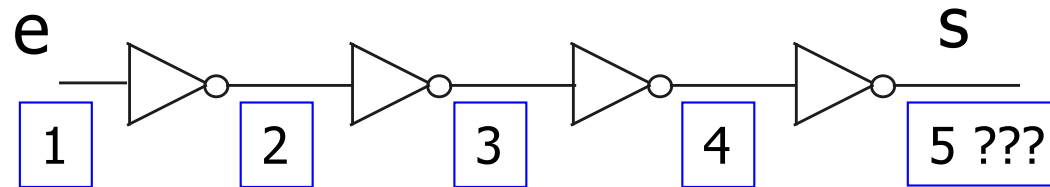
$e_1 = 0$



SCOAP Pros and Cons (2)

SCOAP CC1

SCOAP CC0



However, it is as easy to control s as e



The function structure influences the results



COP Principle

- Based on a structural analysis of the DUT
- COP Measurements → Probabilities
 - Proportion of input vectors that can control and observe a net
- For combinatorial parts of DUT only



COP Measurements

- Five measurements associated to each net:
 - $C1(N)$ = Controllability at 1 of net N
= (# input vectors generating a 1 on N) / 2^n (n = # bits)
 - $C0(N)$ = Controllability at 0 of net N
= (# input vectors generating a 0 on N) / 2^n
= $1 - C1(N)$
 - $O(N)$ = Observability of net N
= (# of 1 and 0 on N observable on at least one PO) / 2^n
 - $O1(N)$ = Observability at 1 of net N
= Detection probability of the Sa0 on net N
= $C1(N).O(N)$
 - $O0(N)$ = Observability at 0 of net N
= Detection probability of the Sa1 on net N
= $C0(N).O(N)$



Classical Logic Gates

- AND

- $C1(s) = C1(e1) \cdot C1(e2)$
- $O(e1) = C1(e2) \cdot O(s)$

- OR

- $C1(s) = 1 - (1 - C1(e1)) \cdot (1 - C1(e2))$
- $O(e1) = (1 - C1(e2)) \cdot O(s)$

- NOT

- $C1(s) = 1 - C1(e)$
- $O(e) = O(s)$



COP Process

- Net initialisation:
 - $C1(\text{PIs}) = 0,5$
 - $O(\text{POs}) = 1$
- Phase 1:
 - Controllability computation form PIs to POs
- Phase 2:
 - Observability computation form POs to PIs



Conclusion

- Controllability and Observability measurements
 - Can represent a “bad” image of the testability
- Possible misinterpretation **BUT** still widely used in industry
 - To help the test vector generation process (decision making)
 - To guide the Design-for-Test steps (test point insertion, ...)